# Degree Spectra and Conservative Extensions of Abstract Structures

Alexandra A. Soskova <sup>1</sup>
joint work with
Ivan N. Soskov and Stefan V. Vatev

Faculty of Mathematics and Computer Science Sofia University

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#### Outline

- Degree spectra of structures
- Definability on structures
- Conservative (k, n) Extensions
- Jumps of Structures
- Jump inversion theorem for structures

## Degree Spectra

Let  $\mathfrak{A} = (A; P_1, \dots, P_k)$  be a denumerable structure. Enumeration of  $\mathfrak{A}$  is every one to one mapping of  $\mathbb{N}$  onto A.

Given an enumeration f of  $\mathfrak A$  and a subset of X of  $A^a$ , let

$$f^{-1}(X) = \{\langle x_1, \dots, x_a \rangle : (f(x_1), \dots, f(x_a)) \in X\}.$$

Set 
$$f^{-1}(\mathfrak{A}) = f^{-1}(P_1) \oplus \cdots \oplus f^{-1}(P_k) \oplus f^{-1}(=) \oplus f^{-1}(\neq)$$
.

**Definition.**(Richter) *The Degree Spectrum of*  $\mathfrak A$  is the set

$$DS_T(\mathfrak{A}) = \{d_T(f^{-1}(\mathfrak{A})) : f \text{ is an enumeration of } \mathfrak{A})\}.$$



#### Degree Spectra

**Definition.**(Knight) The n-th jump spectrum of a structure  $\mathfrak A$  is the set

$$DS_n(\mathfrak{A}) = \{\mathbf{a}^{(n)} | \mathbf{a} \in DS(\mathfrak{A})\}.$$

**Proposition.** (Knight) For every automorphically nontrivial structure  $\mathfrak{A}$ ,  $DS_n(\mathfrak{A})$  is an upwards closed set of degrees.

**Theorem.**(A. Soskova, I. Soskov) Every jump spectrum is a spectrum of a structure, i.e. for every countable structure  $\mathfrak A$  there is a structure  $\mathfrak B$  such that  $DS_1(\mathfrak A) = DS(\mathfrak B)$ .

**Theorem.** (A. Soskova, I. Soskov) Let  $\mathfrak A$  and  $\mathfrak C$  be countable structures and  $DS(\mathfrak A) \subseteq DS_1(\mathfrak C)$ . There exists a structure  $\mathfrak B$  such that  $DS(\mathfrak A) = DS_1(\mathfrak B)$  and  $DS(\mathfrak B) \subseteq DS(\mathfrak C)$ .

## Formally $\Sigma_n^c$ -definable sets

Let L be the language of  $\mathfrak{A}$ . The computable  $\Sigma_n^c$  formulas in L are defined inductively:

- A computable  $\Sigma_0^c$  ( $\Pi_0^c$ ) formula is a finitary quantifier-free formula in L.
- A computable  $\sum_{n=1}^{c}$  formula  $\Phi(\overline{x})$  is a disjunction of c.e. set of formulas of the form

$$(\exists \overline{Y}) \Psi(\overline{X}, \overline{Y})$$

 $\Psi$  is a finite conjunction of  $\Sigma_n^c$  and  $\Pi_n^c$  formulas

•  $\Pi_{n+1}^c$  formulas are the negations of the  $\Sigma_{n+1}^c$  formulas.



## Example

Consider  $\mathcal{O} = (\mathbb{N}; =)$  and  $\mathcal{S} = (\mathbb{N}; G_{Succ}; =)$ , where  $G_{Succ}$  is the graph of the successor function.

$$DS(\mathcal{O}) = DS(\mathcal{S})$$

The  $\Sigma_1^c(\mathcal{O})$  are all finite and co-finite sets of natural numbers. But all c.e. set are formally  $\Sigma_1^c$  definable on  $\mathcal{S}$ . So, the structure  $\mathcal{S}$  is more powerful than the  $\mathcal{O}$ .

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#### **Enumerations**

**Definition.** The pair  $\alpha=(f_{\alpha},R_{\alpha})$  is an enumeration of the set  $X\subseteq A$ , if  $R_{\alpha}$  is a set of natural numbers,  $f_{\alpha}$  is a partial one-to-one mapping of  $\mathbb N$  onto X and  $\mathrm{dom}(f_{\alpha})=f_{\alpha}^{-1}(X)$  is c.e. in  $R_{\alpha}$ . We denote this by  $X\leq \alpha$ .

**Definition.** The pair  $\alpha=(f_{\alpha},R_{\alpha})$  is an *enumeration* of  $\mathfrak A$  if  $\alpha$  is an enumeration of A and  $f_{\alpha}^{-1}(\mathfrak A)$  is computable in  $R_{\alpha}$ . We denote this by  $\mathfrak A \leq \alpha$ .

For an enumeration  $\alpha = (f_{\alpha}, R_{\alpha})$  of  $\mathfrak{A}$  we denote by  $\alpha^{(n)} = (f_{\alpha}, R_{\alpha}^{(n)}).$ 



#### Reformulation

The Degree Spectrum of  $\mathfrak A$  is the set

$$DS(\mathfrak{A}) = \{ d_T(R_\alpha) \mid \mathfrak{A} \leq \alpha \}.$$

**Theorem.** (Ash, Knigh, Manasse, Slaman, Chisholm) For every set  $X \subseteq A$ ,

$$X \in \Sigma_{n+1}^{c}(\mathfrak{A}) \leftrightarrow (\forall \alpha)[\mathfrak{A} \leq \alpha \to X \leq \alpha^{(n)}].$$

#### Conservative (k, n) Extensions

Let  $\alpha = (f_{\alpha}, R_{\alpha})$  and  $\beta = (f_{\beta}, R_{\beta})$  be enumerations of the structures  $\mathfrak{A}$  and  $\mathfrak{B}$  respectively. We write  $\alpha < \beta$  if

- (i)  $R_{\alpha} \leq_{T} R_{\beta}$  and
- (ii) the set  $E(f_{\alpha}, f_{\beta}) = \{(x, y) \mid x \in Dom(f_{\alpha}) \& y \in Dom(f_{\beta}) \& f_{\alpha}(x) = f_{\beta}(y)\}$  is c.e. in  $R_{\beta}$ .

## Conservative (k, n) Extensions

**Definition.** Let  $\mathfrak A$  and  $\mathfrak B$  be countable structures, possibly with different signatures and  $A\subseteq B$ .

- (i)  $\mathfrak{A} \leq_n^k \mathfrak{B}$  iff for every enumeration  $\beta$  of  $\mathfrak{B}$  there exists an enumeration  $\alpha$  of  $\mathfrak{A}$  such that  $\alpha^{(k)} \leq \beta^{(n)}$ .
- (ii)  $\mathfrak{A} \geq_n^k \mathfrak{B}$  iff for every enumeration  $\alpha$  of  $\mathfrak{A}$  there exists an enumeration  $\beta$  of  $\mathfrak{B}$  such that  $\beta^{(n)} \leq \alpha^{(k)}$ .
- (iii)  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  if  $\mathfrak{A} \leq_n^k \mathfrak{B}$  and  $\mathfrak{A} \geq_n^k \mathfrak{B}$ . We shall say that  $\mathfrak{B}$  is a (k, n)-conservative extension of  $\mathfrak{A}$ .

Note that the relation  $\equiv_n^k$  is not symmetric.



# Conservative (k, n) Extensions and Degree Spectra

**Proposition.** Let  $\mathfrak A$  and  $\mathfrak B$  be countable structures with  $A\subseteq B$ .

- (i) If  $\mathfrak{A} \leq_n^k \mathfrak{B}$  then  $DS_n(\mathfrak{B}) \subseteq DS_k(\mathfrak{A})$ ;
- (ii) If  $\mathfrak{A} \geq_n^k \mathfrak{B}$  then  $DS_k(\mathfrak{A}) \subseteq DS_n(\mathfrak{B})$ ;
- (iii) If  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  then  $DS_k(\mathfrak{A}) = DS_n(\mathfrak{B})$ ;

#### Corollary.

- (i) k = 1, n = 0: If  $\mathfrak{A} \equiv_0^1 \mathfrak{B}$  then  $DS_1(\mathfrak{A}) = DS(\mathfrak{B})$ .
- (ii) k = 0, n = 1: If  $\mathfrak{A} \equiv_1^0 \mathfrak{B}$  then  $DS(\mathfrak{A}) = DS_1(\mathfrak{B})$ .



## Conservative (k, n) Extensions and Definability

**Theorem.** Let for  $\mathfrak A$  and  $\mathfrak B$ :  $A\subseteq B$ . For all  $k,n\in\mathbb N$ ,

- (i) if  $\mathfrak{A} \leq_n^k \mathfrak{B}$  then  $(\forall X \subseteq A)[X \in \Sigma_{k+1}^c(\mathfrak{A}) \to X \in \Sigma_{n+1}^c(\mathfrak{B})]$ ;
- (ii) if  $\mathfrak{A} \geq_n^k \mathfrak{B}$  then  $(\forall X \subseteq A)[X \in \Sigma_{n+1}^c(\mathfrak{B}) \to X \in \Sigma_{k+1}^c(\mathfrak{A})]$ ;
- (iii) if  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  then  $(\forall X \subseteq A)[X \in \Sigma_{k+1}^c(\mathfrak{A}) \leftrightarrow X \in \Sigma_{n+1}^c(\mathfrak{B})]$ .

## Conservative (k, n) Extensions and Definability

The opposite direction is not always true:

#### Example.

Consider  $\mathcal{O}_A = (A; =)$  and take  $\mathfrak{A} = \mathfrak{B} = \mathcal{O}_A$ .

For every natural number n,

 $X \subseteq A$  is  $\Sigma_n^c(\mathcal{O}_A)$  iff X is a finite or co-finite subset of A.

Therefore  $\Sigma_1^c(\mathcal{O}_A) = \Sigma_n^c(\mathcal{O}_A)$  and

$$(\forall n)(\forall X \subseteq A)[X \in \Sigma_{n+1}^c(\mathcal{O}_A) \to X \in \Sigma_1^c(\mathcal{O}_A)].$$

But  $(\forall n)[\mathcal{O}_A \leq_0^n \mathcal{O}_A]$  is evidently not true.

#### Moschovakis' extension

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Let \mathfrak{A}=(A;P_1,\ldots,P_k) and \bar{0}\not\in A.

Set A_0=A\cup\{\bar{0}\}.

Let \langle .,.\rangle be a pairing function s.t. none of the elements of A is a pair and A^* be the least set containing A_0 and closed under \langle .,.\rangle.

Let 0^*=\bar{0} and (n+1)^*=\langle \bar{0},n^*\rangle,\ \mathbb{N}^*=\{n^*\mid n\in\mathbb{N}\}.

The decoding functions: L(\langle s,t\rangle)=s\ \&\ R(\langle s,t\rangle)=t

L(\bar{0})=R(\bar{0})=0^*\ (\forall t\in A)[L(t)=R(t)=1^*].
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#### Moschovakis' extension

**Definition.** *Moschovakis' extension* of  $\mathfrak A$  is the structure

$$\mathfrak{A}^{\star} = (A^{\star}, P_1, \dots, P_k, A_0, G_{\langle \dots \rangle}, G_L, G_R).$$

**Proposition.**  $\mathfrak{A} \equiv_n^n \mathfrak{A}^*$  for every  $n \in \mathbb{N}$ .

**Proposition.** For every two structures  $\mathfrak{A}$ ,  $\mathfrak{B}$  with  $A \subseteq B$  and natural numbers n, k  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  iff  $\mathfrak{A}^* \equiv_n^k \mathfrak{B}^*$ .

# Conservative (k, n) Extensions and Definability

#### **Theorem.**(S. Vatev)

Let  $\mathfrak A$  and  $\mathfrak B$  be countable structures with  $A^\star\subseteq B$  and  $k,n\in\mathbb N$ . If  $(\forall X\subseteq A^\star)[X\in\Sigma_{k+1}^c(\mathfrak A^\star)\to X\in\Sigma_{n+1}^c(\mathfrak B)]$  then  $\mathfrak A\le_n^k\mathfrak B$ .

**Corollary.** For any two countable structures  $\mathfrak{A}$ ,  $\mathfrak{B}$  with  $A \subseteq B$  and  $n, k \in \mathbb{N}$ ,

$$\mathfrak{A} \leq_n^k \mathfrak{B} \leftrightarrow (\forall X \subseteq A^*)[X \in \Sigma_{k+1}^c(\mathfrak{A}^*) \to X \in \Sigma_{n+1}^c(\mathfrak{B}^*)].$$

#### The set $K^{\mathfrak{A}}$

A new predicate  $K_{\mathfrak{A}}$  (analogue of Kleene's set). For  $e, x \in \mathbb{N}$  and finite part  $\tau$ , let

$$\tau \Vdash F_e(x) \leftrightarrow x \in W_e^{\tau^{-1}(\mathfrak{A})}$$

$$\tau \Vdash \neg F_e(x) \leftrightarrow (\forall \rho \supseteq \tau)(\rho \not\Vdash F_e(x))$$

$$K^{\mathfrak{A}} = \{ \langle \delta, e, x \rangle : (\exists \tau \supseteq \delta)(\tau \Vdash F_e(x)) \}.$$

$$\mathfrak{A}' = (\mathfrak{A}^*, K^{\mathfrak{A}}).$$

Theorem.  $DS_1(\mathfrak{A}) = DS(\mathfrak{A}')$ .

**Proposition.**  $\mathfrak{A} \equiv_0^1 \mathfrak{A}'$ .



## The Forcing Relation

For every  $e, x, n \in \mathbb{N}$  and for every finite part  $\tau : \mathbb{N} \to A$ , we define the forcing relations  $\Vdash_n$ :

$$\tau \Vdash_{0} F_{e}(x) \qquad \leftrightarrow \qquad x \in W_{e}^{\tau^{-1}(\mathfrak{A})}$$

$$\tau \Vdash_{n+1} F_{e}(x) \qquad \leftrightarrow \qquad (\exists v)[\langle x, v \rangle \in W_{e} \& (\forall u \in D_{v})[(u = \langle e_{u}, x_{u}, 1 \rangle \& \tau \Vdash_{n} F_{e_{u}}(x_{u})) \lor (u = \langle e_{u}, x_{u}, 0 \rangle \& \tau \Vdash_{n} \neg F_{e_{u}}(x_{u}))]],$$

$$\tau \Vdash_{n} \neg F_{e}(x) \qquad \leftrightarrow \qquad (\forall \rho \supseteq \tau)(\rho \not\Vdash_{n} F_{e}(x)).$$

# The set $K_n^{\mathfrak{A}}$

#### Definition.

$$K_n^{\mathfrak{A}} = \{ \langle \delta, e, x \rangle \mid (\exists \tau \supseteq \delta) [\tau \Vdash_n F_e(x)] \}.$$

#### **Proposition.**(S. Vatev)

- (i)  $K_n^{\mathfrak{A}} \in \Sigma_{n+1}^c(\mathfrak{A}^\star)$  and  $A^\star \setminus K_n^{\mathfrak{A}} \in \Sigma_{n+2}^c(\mathfrak{A}^\star)$ .
- (ii)  $K_n^{\mathfrak{A}} \notin \Sigma_n^c(\mathfrak{A}^*)$ .

#### Jumps of Structures

**Definition.** For every natural number n, we define the n-th jump of the structure  $\mathfrak A$  in the following way:

$$\mathfrak{A}^{(0)} = \mathfrak{A}$$
 and  $\mathfrak{A}^{(n+1)} = (\mathfrak{A}^{\star}, K_n^{\mathfrak{A}}).$ 

## Jumps of Structures

**Proposition.** For every  $\mathfrak A$  and natural number n,

- (i)  $\mathfrak{A} \equiv_0^n \mathfrak{A}^{(n)}$ ;
- (ii)  $\mathfrak{A}^{(n)} \leq_0^0 \mathfrak{A}^{(n+1)}$  and  $\mathfrak{A}^{(n)} \not\equiv_0^0 \mathfrak{A}^{(n+1)}$ .

Since  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  implies  $DS_k(\mathfrak{A}) = DS_n(\mathfrak{B})$ , we get the following.

**Corollary.** For every  $\mathfrak{A}$ ,  $DS(\mathfrak{A}^{(n)}) = DS_n(\mathfrak{A})$ .

## The Jump Inversion Theorem

**Theorem.** Let  $\mathfrak{A}$  and  $\mathfrak{C}$  be countable structures and  $DS(\mathfrak{A}) \subseteq DS_1(\mathfrak{C})$ . There exists a structure  $\mathfrak{B} = \mathfrak{A}^{\exists \forall} \oplus \mathfrak{C}$  such that  $DS(\mathfrak{A}) = DS_1(\mathfrak{B})$  and  $DS(\mathfrak{B}) \subseteq DS(\mathfrak{C})$ .

#### Remark. Similar results by:

- A. Montalban (2009) by different approach with complete set of  $\Pi_n^c$  formulas.
- A. Stukachev (2009) for  $\Sigma$  reducibility with Marker's extentions.

Stukachev proves an analogue of this theorem for the semilattices of  $\Sigma$ -degrees of structures with arbitrary cardinalities.

**Theorem.**(Stukachev) Let  $\mathfrak A$  be a structure such that  $\mathbf 0' \leq_{\mathbf \Sigma} \mathfrak A$ . There exists a structure  $\mathfrak B$  such that  $\mathfrak A \equiv_{\mathbf \Sigma} \mathfrak B'$ .

We can prove a similar to Stukachev's result.

## The Jump Inversion Theorem

**Proposition.** If  $\mathcal{O}_A \leq_0^1 \mathfrak{A}$ , then  $\mathfrak{A} \equiv_1^0 \mathfrak{A}^{\exists \forall}$ .

**Theorem.** Let  $\mathcal{O}_A \leq_0^k \mathfrak{A}$  for some  $k \in \mathbb{N}$ . There exists a structure  $\mathfrak{B} = \mathfrak{A}^{\exists \forall}$  such that  $\mathfrak{A} \equiv_0^0 \mathfrak{B}^{(k)}$ .

**Remark.** Note that  $\mathcal{O}_A \leq_0^k \mathfrak{A}$  iff the elements of  $DS(\mathfrak{A})$  are above  $\mathbf{0}^{(k)}$ .

## The Jump Inversion Theorem

**Proposition.** Let  $\mathcal{O}_A \leq_0^k \mathfrak{A}$  for some  $k \in \mathbb{N}$ . There exists a structure  $\mathfrak{B}$  such that for every  $n \in \mathbb{N}$ ,  $\mathfrak{A} \equiv_k^n \mathfrak{B}^{(n)}$ .

**Corollary.** Let  $\mathcal{O}_A \leq_0^k \mathfrak{A}$  for some  $k \in \mathbb{N}$ . There exists a countable structure  $\mathfrak{B}$  such that

$$(\forall n \in \mathbb{N})(\forall X \subseteq A)[X \in \Sigma_{n+1}^{c}(\mathfrak{A}) \leftrightarrow X \in \Sigma_{k+1}^{c}(\mathfrak{B}^{(n)})].$$

**Corollary.** If  $\mathcal{O}_A \leq_0^k \mathfrak{A}$  for some  $k \in \mathbb{N}$  then for each  $n \in \mathbb{N}$ , there is a structure  $\mathfrak{B}$  such that

$$(\forall X\subseteq A)[X\in \Sigma_{n+1}^c(\mathfrak{A})\leftrightarrow X\in \Sigma_{k+1}^c(\mathfrak{B})].$$



#### Some problems

- The definition of  $\mathfrak{A} \equiv_n^k \mathfrak{B}$  is not symmetric since we suppose that  $A \subseteq B$ . How to define the similar relation more symmetric and for arbitrary  $\mathfrak{A}$  and  $\mathfrak{B}$ ?
- How to relativize the Jump Inversion Theorem for structures?
- The Jump inversion Theorem for structures for arbitrary constructive ordinal  $\alpha$ .

## Thank you!

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