

Properties of degree spectra , co-spectra and omega co-spectra

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Enumeration reducibility

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$A \leq_e B$ if there is a c.e. set W , such that

$$A = W(B) = \{x \mid \exists D(\langle x, D \rangle \in W \ \& \ D \subseteq B)\}.$$

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Selman's theorem

Equivalently, $A \leq_e B$ if there is a single Turing functional which uniformly, given any enumeration of B , outputs an enumeration of A .

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Theorem (Selman)

The set A is enumeration reducible to the set B if and only if $\mathcal{E}(B) \subseteq \mathcal{E}(A)$.

The enumeration jump

Definition. Given a set A , denote by $A^+ = A \oplus (\mathbb{N} \setminus A)$.

Theorem. For any sets A and B :

- 1 A is c.e. in B iff $A \leq_e B^+$.
- 2 $A \leq_T B$ iff $A^+ \leq_e B^+$.

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Definition. For any set A let $K_A = \{\langle i, x \rangle \mid x \in W_i(A)\}$. Set $A' = K_A^+$.

Definition. A set A is called *total* iff $A \equiv_e A^+$.

Let $d_e(A)' = d_e(A')$. The enumeration jump is always a total degree and agrees with the Turing jump under the standard embedding $\iota : \mathcal{D}_T \rightarrow \mathcal{D}_e$ by $\iota(d_T(A)) = d_e(A^+)$.

Enumeration degree spectra

Let $\mathfrak{A} = (A; R_1, \dots, R_k)$ be a countable structure. An enumeration of \mathfrak{A} is every total surjective mapping of \mathbb{N} onto A .

Given an enumeration f of \mathfrak{A} and a subset B of A^n , let

$$f^{-1}(B) = \{\langle x_1, \dots, x_n \rangle \mid (f(x_1), \dots, f(x_n)) \in B\}.$$

$$f^{-1}(\mathfrak{A}) = f^{-1}(R_1) \oplus \dots \oplus f^{-1}(R_k) \oplus f^{-1}(=) \oplus f^{-1}(\neq).$$

Definition. The enumeration degree spectrum of \mathfrak{A} is the set

$$DS(\mathfrak{A}) = \{d_e(f^{-1}(\mathfrak{A})) \mid f \text{ is an enumeration of } \mathfrak{A}\}.$$

If \mathbf{a} is the least element of $DS(\mathfrak{A})$, then \mathbf{a} is called the *e-degree* of \mathfrak{A} .

Enumeration degree spectra

Proposition. *The enumeration degree spectrum is closed upwards with respect to total e-degrees, i.e. if $\mathbf{a} \in DS(\mathfrak{A})$, \mathbf{b} is a total e-degree $\mathbf{a} \leq_e \mathbf{b}$ then $\mathbf{b} \in DS(\mathfrak{A})$.*

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Let $\mathfrak{A}^+ = (A, R_1, \dots, R_k, R_1^c, \dots, R_k^c)$.

Proposition.

$\iota(DS_T(\mathfrak{A})) = DS(\mathfrak{A}^+)$.

Co-spectra

Definition. Let \mathcal{A} be a nonempty set of enumeration degrees. The *co-set* of \mathcal{A} is the set $co(\mathcal{A})$ of all lower bounds of \mathcal{A} . Namely

$$co(\mathcal{A}) = \{\mathbf{b} : \mathbf{b} \in \mathcal{D}_e \ \& \ (\forall \mathbf{a} \in \mathcal{A})(\mathbf{b} \leq_e \mathbf{a})\}.$$

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Definition. Given a structure \mathfrak{A} , set $CS(\mathfrak{A}) = co(DS(\mathfrak{A}))$.
If \mathbf{a} is the greatest element of $CS(\mathfrak{A})$ then we call \mathbf{a} the *co-degree* of \mathfrak{A} .

The admissible in \mathfrak{A} sets

Definition. A set B of natural numbers is admissible in \mathfrak{A} if for every enumeration f of \mathfrak{A} , $B \leq_e f^{-1}(\mathfrak{A})$.

Clearly $\mathbf{a} \in CS(\mathfrak{A})$ iff $\mathbf{a} = d_e(B)$ for some admissible in \mathfrak{A} set B .

The formally definable sets on \mathfrak{A}

Definition. A Σ_1^+ formula with free variables among X_1, \dots, X_r is a c.e. disjunction of existential formulae of the form $\exists Y_1 \dots \exists Y_k \theta(\vec{Y}, \vec{X})$, where θ is a finite conjunction of atomic formulae.

Definition. A set $B \subseteq \mathbb{N}$ is *formally definable* on \mathfrak{A} if there exists a computable function $\gamma(x)$, such that $\bigvee_{x \in \mathbb{N}} \Phi_{\gamma(x)}$ is a Σ_1^+ formula with free variables among X_1, \dots, X_r and elements t_1, \dots, t_r of A such that the following equivalence holds:

$$x \in B \iff \mathfrak{A} \models \Phi_{\gamma(x)}(X_1/t_1, \dots, X_r/t_r) .$$

Theorem. Let $B \subseteq \mathbb{N}$. Then

- 1 $d_e(B) \in CS(\mathfrak{A})$ iff
- 2 $B \leq_e f^{-1}(\mathfrak{A})$ for all generic enumerations f of \mathfrak{A} iff
- 3 B is formally definable on \mathfrak{A} .

Jump spectra and jump co-spectra

Definition. The n th jump spectrum of \mathfrak{A} is the set

$$DS_n(\mathfrak{A}) = \{d_e(f^{-1}(\mathfrak{A})^{(n)}) : f \text{ is an enumeration of } \mathfrak{A}\}.$$

If \mathbf{a} is the least element of $DS_n(\mathfrak{A})$, then \mathbf{a} is called the n th jump degree of \mathfrak{A} .

Definition. The co-set $CS_n(\mathfrak{A})$ of the n th jump spectrum of \mathfrak{A} is called n th jump co-spectrum of \mathfrak{A} .

If $CS_n(\mathfrak{A})$ has a greatest element then it is called the n th jump co-degree of \mathfrak{A} .

Some examples

- For every linear ordering \mathfrak{A} $DS(\mathfrak{A})$ contains a minimal pair of degrees [Richter] and hence $\mathbf{0}_e$ is the co-degree of \mathfrak{A} . So, if \mathfrak{A} has a degree \mathbf{a} , then $\mathbf{a} = \mathbf{0}_e$.
- For a linear ordering \mathfrak{A} , $CS_1(\mathfrak{A})$ consists of all e-degrees of Σ_2^0 sets [Knight]. The first co-degree of \mathfrak{A} is $\mathbf{0}'_e$.
- There exists a structure \mathfrak{A} [Slaman, Whener]

$$DS(\mathfrak{A}) = \{\mathbf{a} : \mathbf{a} \text{ is total and } \mathbf{0}_e < \mathbf{a}\}.$$

Clearly, the structure \mathfrak{A} has co-degree $\mathbf{0}_e$ but has no degree.

- There is a structure whose spectrum is exactly the non-hyperarithmetical degrees [Greenberg, Motalbán and Slaman]

A special kind of co-degree

Definition. [Knight, Motalbán] A structure \mathfrak{A} has “enumeration degree X ” if every enumeration of X computes a copy of \mathfrak{A} , and every copy of \mathfrak{A} computes an enumeration of X .

In our terms this can be formulated as \mathfrak{A}^+ has a co-degree $d_e(X)$ and $DS(\mathfrak{A}) = \{\mathbf{a} \mid \mathbf{a} \text{ is total and } d_e(X) \leq \mathbf{a}\}$.

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Example. Given $X \subseteq \mathbb{N}$, consider the group $G_X = \bigoplus_{i \in X} \mathbb{Z}_{p_i}$, where p_i is the i th prime number. Then G_X has “enumeration degree X ”: We can easily build G_X out of an enumeration of X , and for the other direction, we have that $n \in X$ if and only if there exists $g \in G_X$ of order p_n .

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Theorem. [A. Montalbán] Let K be Π_2^c class of \exists -atomic structures, i.e. K is the class of structures axiomatized by some Π_2^c sentence and for every structure \mathfrak{A} in K and every tuple $\bar{a} \in |\mathfrak{A}|$ the orbit of \bar{a} is existentially definable (with parameters \bar{a}). Then every structure in K has “enumeration degree” given by its \exists -theory.

Representing the principle countable ideals as co-spectra

Example. Let G be a torsion free abelian group of rank 1. [Coles, Downey, Slaman; Soskov] There exists an enumeration degree \mathbf{s}_G such that

- $DS(G) = \{\mathbf{b} : \mathbf{b} \text{ is total and } \mathbf{s}_G \leq \mathbf{b}\}$.
- The co-degree of G is \mathbf{s}_G .
- G has a degree iff \mathbf{s}_G is a total e-degree.

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- G has a degree iff \mathbf{s}_G is a total e-degree.

For every $\mathbf{d} \in \mathcal{D}_e$ there exists a G , s.t. $\mathbf{s}_G = \mathbf{d}$.

Corollary. Every principle ideal of enumeration degrees is $CS(G)$ for some G .

Representing non-principle countable ideals as co-spectra

Theorem. [Soskov] *Every countable ideal is the co-spectrum of a structure.*

Proof.

Let B_0, \dots, B_n, \dots be a sequence of sets of natural numbers. Set $\mathfrak{A} = (\mathbb{N}; \mathbf{G}_f; \sigma)$,

$$f(\langle i, n \rangle) = \langle i + 1, n \rangle;$$

$$\sigma = \{ \langle i, n \rangle : n = 2k + 1 \vee n = 2k \ \& \ i \in B_k \}.$$

Then $CS(\mathfrak{A}) = I(d_e(B_0), \dots, d_e(B_n), \dots)$



Spectra with a countable base

Definition. Let $\mathcal{B} \subseteq \mathcal{A}$ be sets of degrees. Then \mathcal{B} is a base of \mathcal{A} if

$$(\forall \mathbf{a} \in \mathcal{A})(\exists \mathbf{b} \in \mathcal{B})(\mathbf{b} \leq \mathbf{a}).$$

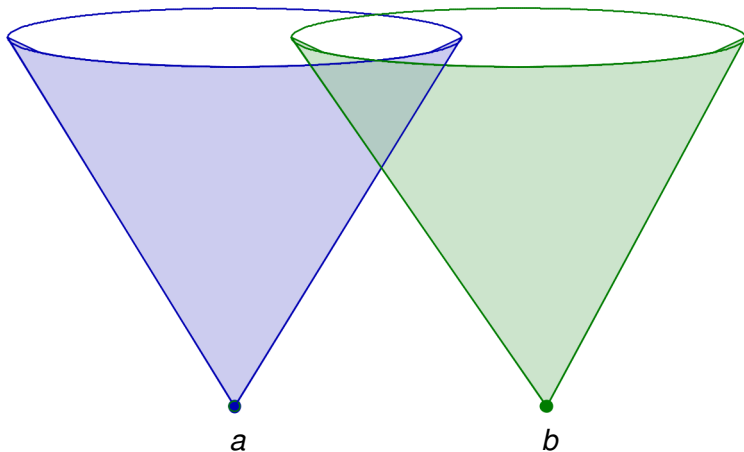
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Theorem. *A structure \mathfrak{A} has e-degree if and only if $DS(\mathfrak{A})$ has a countable base.*

An upwards closed set of degrees which is not a degree spectra of a structure



Other examples

- The class of PA degrees is not the degree spectrum of any structure [[Andrews, Miller](#)].
- The upward closure of the set of 1-random degrees is not the spectrum of a structure [[Andrews, Miller](#)].
- A degree spectrum is never the Turing-upward closure of F_σ set of reals in ω^ω , unless it is enumeration cone [[Montalbà](#)]

The minimal pair theorem

Theorem. Let $\mathbf{c} \in DS_n(\mathfrak{A})$. There exist total $\mathbf{f}, \mathbf{g} \in DS(\mathfrak{A})$ such that, $\mathbf{f}^{(n)} = \mathbf{g}^{(n)} \leq \mathbf{c}$ and $CS_k(\mathfrak{A}) = co(\{\mathbf{f}^{(k)}, \mathbf{g}^{(k)}\})$ for every $k \leq n - 2$.

Notice that for every enumeration degree \mathbf{b} there exists a structure $\mathfrak{A}_{\mathbf{b}}$ such that $DS(\mathfrak{A}_{\mathbf{b}}) = \{\mathbf{x} \in \mathcal{D}_T \mid \mathbf{b} <_e \mathbf{x}\}$. Hence

Corollary. [Rozinas] For every $\mathbf{b} \in \mathcal{D}_e$ there exist total \mathbf{f}, \mathbf{g} below \mathbf{b}'' which are a minimal pair over \mathbf{b} .

The quasi-minimal degree

Definition. Let \mathcal{A} be a set of enumeration degrees. The degree \mathbf{q} is quasi-minimal with respect to \mathcal{A} if:

- $\mathbf{q} \notin \text{co}(\mathcal{A})$.
- If \mathbf{a} is total and $\mathbf{a} \geq \mathbf{q}$, then $\mathbf{a} \in \mathcal{A}$.
- If \mathbf{a} is total and $\mathbf{a} \leq \mathbf{q}$, then $\mathbf{a} \in \text{co}(\mathcal{A})$.

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Theorem. *For every structure \mathfrak{A} there exists a quasi-minimal with respect to $DS(\mathfrak{A})$ degree.*

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Theorem. For every structure \mathfrak{A} there exists a quasi-minimal with respect to $DS(\mathfrak{A})$ degree.

Corollary. [Slaman and Sorbi] Let I be a countable ideal of enumeration degrees. There exists an enumeration degree \mathbf{q} s.t.

- 1 If $\mathbf{a} \in I$ then $\mathbf{a} <_e \mathbf{q}$.
- 2 If \mathbf{a} is total and $\mathbf{a} \leq_e \mathbf{q}$ then $\mathbf{a} \in I$.

Jumps of quasi-minimal degrees

Proposition. *For every countable structure \mathfrak{A} there exist uncountably many quasi-minimal degrees with respect to $DS(\mathfrak{A})$.*

Proposition. *The first jump spectrum of every structure \mathfrak{A} consists exactly of the enumeration jumps of the quasi-minimal degrees.*

Corollary. [McEvoy] *For every total e -degree $\mathbf{a} \geq_e \mathbf{0}'_e$ there is a quasi-minimal degree \mathbf{q} with $\mathbf{q}' = \mathbf{a}$.*

Splitting a total set

Proposition. [*Jockusch*] For every total e-degree \mathbf{a} there are quasi-minimal degrees \mathbf{p} and \mathbf{q} such that $\mathbf{a} = \mathbf{p} \vee \mathbf{q}$.

Proposition. For every element \mathbf{a} of the jump spectrum of a structure \mathfrak{A} there exists quasi-minimal with respect to $DS(\mathfrak{A})$ degrees \mathbf{p} and \mathbf{q} such that $\mathbf{a} = \mathbf{p} \vee \mathbf{q}$.

Every jump spectrum is the spectrum of a structure

Let $\mathfrak{A} = (A; R_1, \dots, R_n)$.

Let $\bar{0} \notin A$. Set $A_0 = A \cup \{\bar{0}\}$. Let $\langle \cdot, \cdot \rangle$ be a pairing function s.t. none of the elements of A_0 is a pair and A^* be the least set containing A_0 and closed under $\langle \cdot, \cdot \rangle$. Let L and R be the decoding functions.

Definition. *Moschovakis'* extension of \mathfrak{A} is the structure

$$\mathfrak{A}^* = (A^*, R_1, \dots, R_n, A_0, G_{\langle \cdot, \cdot \rangle}, G_L, G_R).$$

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Let $K_{\mathfrak{A}} = \{ \langle \delta, e, x \rangle : (\exists \tau \supseteq \delta)(\tau \Vdash F_e(x)) \}$.

Set $\mathfrak{A}' = (\mathfrak{A}^*, K_{\mathfrak{A}}, A^* \setminus K_{\mathfrak{A}})$.

Theorem. $DS_1(\mathfrak{A}) = DS(\mathfrak{A}')$.

The jump inversion theorem

Let $\alpha < \omega_1^{CK}$ and \mathfrak{A} be a countable structure such that all elements of $DS(\mathfrak{A})$ are above $\mathbf{0}^{(\alpha)}$.

Does there exist a structure \mathfrak{M} such that $DS_\alpha(\mathfrak{M}) = DS(\mathfrak{A})$?

Theorem. [Soskov, AS] $\alpha = 1$. If $DS(\mathfrak{A}) \subseteq DS_1(\mathfrak{B})$ then there exists a structure \mathfrak{C} such that $DS(\mathfrak{C}) \subseteq DS(\mathfrak{B})$ and $DS_1(\mathfrak{C}) = DS(\mathfrak{A})$.

Method: Marker's extensions.

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Method: Marker's extensions.

Remark. If a structure \mathfrak{B} has the property $DS(\mathfrak{A}) = DS_1(\mathfrak{B})$, then it follows that $DS(\mathfrak{C}) \subseteq DS(\mathfrak{B})$.

2009 [Montalbà](#) Notes on the jump of a structure, *Mathematical Theory and Computational Practice*, 372–378.

2009 [Stukachev](#) A jump inversion theorem for the semilattices of Sigma-degrees, *Siberian Electronic Mathematical Reports*, v. 6, 182 – 190

Jump inversion theorem for ordinals

- The jump inversion theorem holds for successor ordinals
[Goncharov-Harizanov-Knight-McCoy-Miller-Solomon, 2006;
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[Goncharov-Harizanov-Knight-McCoy-Miller-Solomon, 2006;
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- The jump inversion theorem does not hold for $\alpha = \omega$.
[Soskov 2013]

Every member of $\mathbf{a} \in CS_\omega(\mathfrak{N})$ is bounded by a total degree \mathbf{b} , which is also a member of $CS_\omega(\mathfrak{N})$.

Strong jump inversion

If a set Y computes a copy of \mathfrak{A}' then its degree is in $DS_1(\mathfrak{A})$ since $DS_1(\mathfrak{A}) = DS(\mathfrak{A}')$.

This means that there is a set X such that $X' \equiv Y$ and the degree of X computes a copy of \mathfrak{A} , i.e. it is in $DS(\mathfrak{A})$.

Definition. A structure \mathfrak{A} admits a strong jump inversion if for every set X if X' computes a copy of \mathfrak{A}' then X computes a copy of \mathfrak{A} . Equivalently, if \mathfrak{A} has a copy low over X , (the atomic diagram of the copy), then \mathfrak{A} has a computable in X copy.

Proposition. Let $DS(\mathfrak{A}) = DS_1(\mathfrak{B})$.

- *There exists a structure \mathfrak{C} such that $DS(\mathfrak{C}) \subseteq DS(\mathfrak{B})$ and $DS_1(\mathfrak{C}) = DS(\mathfrak{A})$ (by JIT)*
- *If \mathfrak{B} admits a strong jump inversion then for every structure \mathfrak{D} with $DS_1(\mathfrak{D}) = DS(\mathfrak{A}) \implies DS(\mathfrak{D}) \subseteq DS(\mathfrak{B})$.*

Strong jump inversion

- Every Boolean algebra admits strong jump inversion [[Downey and Jockusch](#)]
- There are linear orderings with no computable copy [[Jockusch and Soare](#)]

Strong jump inversion

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- There are linear orderings with no computable copy [[Jockusch and Soare](#)]
- Some sufficient model theoretic conditions, expressed in terms of saturation and enumeration properties of sets of types with formulas of low complexity which guarantee strong jump inversion: [[Calvert, Frolov, Harizanov, Knight, McCoy, AS and Vatev](#)]
- Linear orderings with bounded size of the maximal discrete chains and each element lies in such a chain;
- Linear orderings \mathfrak{A} for which the quotient \mathfrak{A}/\sim is dense and every infinite interval has arbitrary large finite successor chains;
- Abelian p groups of length ω such that the divisible part has infinite dimension;
- Equivalence structures with infinitely many infinite classes;
- Some special trees.

ω -Enumeration Degrees

- Uniform reducibility on sequences of sets.
- For the sequence of sets of natural numbers $\mathcal{B} = \{B_n\}_{n < \omega}$ call *the jump class of \mathcal{B}* the set

$$J_{\mathcal{B}} = \{d_T(X) \mid (\forall n)(B_n \text{ is c.e. in } X^{(n)} \text{ uniformly in } n)\} .$$

Definition. $\mathcal{A} \leq_{\omega} \mathcal{B}$ (\mathcal{A} is ω -enumeration reducible to \mathcal{B}) if $J_{\mathcal{B}} \subseteq J_{\mathcal{A}}$

- $\mathcal{A} \equiv_{\omega} \mathcal{B}$ if $J_{\mathcal{A}} = J_{\mathcal{B}}$.

ω -Enumeration Degrees

- The relation \leq_ω induces a partial ordering of \mathcal{D}_ω with least element $\mathbf{0}_\omega = d_\omega(\emptyset_\omega)$, where \emptyset_ω is the sequence with all members equal to \emptyset .
- \mathcal{D}_ω is further an upper semi-lattice, with least upper bound induced by $\mathcal{A} \oplus \mathcal{B} = \{A_n \oplus B_n\}_{n < \omega}$.
- If $A \subseteq \mathbb{N}$ denote by $A \uparrow \omega = \{A, \emptyset, \emptyset, \dots\}$.
- The mapping $\kappa(d_e(A)) = d_\omega(A \uparrow \omega)$ gives an isomorphic embedding of \mathcal{D}_e to \mathcal{D}_ω , where $A \uparrow \omega = \{A, \emptyset, \emptyset, \dots\}$.

ω -Enumeration Degrees

Let $\mathcal{B} = \{B_n\}_{n < \omega}$.

The jump sequence $\mathcal{P}(\mathcal{B}) = \{\mathcal{P}_n(\mathcal{B})\}_{n < \omega}$:

1 $\mathcal{P}_0(\mathcal{B}) = B_0$

2 $\mathcal{P}_{n+1}(\mathcal{B}) = (\mathcal{P}_n(\mathcal{B}))' \oplus B_{n+1}$

Definition. \mathcal{A} is enumeration reducible \mathcal{B} ($\mathcal{A} \leq_e \mathcal{B}$) iff $A_n \leq_e B_n$ uniformly in n .

Theorem. [Soskov, Kovachev] $\mathcal{A} \leq_\omega \mathcal{B} \iff \mathcal{A} \leq_e \mathcal{P}(\mathcal{B})$.

ω -Enumeration Jump

Definition. The ω -enumeration jump of \mathcal{A} is $\mathcal{A}' = \{\mathcal{P}_{n+1}(\mathcal{A})\}_{n < \omega}$

- $J'_{\mathcal{A}} = \{\mathbf{a}' \mid \mathbf{a} \in J_{\mathcal{A}}\}$.
- The jump is monotone and agrees with the enumeration jump.
- **Soskov and Ganchev:** Strong jump inversion theorem: for $\mathbf{a}^{(n)} \leq \mathbf{b}$ there exists a *least* $\mathbf{x} \geq \mathbf{a}$ such that $\mathbf{x}^{(n)} = \mathbf{b}$. So, every degree \mathbf{x} in the range of the jump operator has a least jump invert.
- **Soskov and Ganchev:** if we add a predicate for the jump operator to the language of partial orders then the natural copy of the enumeration degrees in the omega enumeration degrees becomes first order definable.
- The two structures have the same automorphism group.
- **Ganchev and Sariev:** The jump operator in the upper semi-lattice of the ω -enumeration degrees is first order definable.

ω - Degree Spectra

Let $\mathfrak{A} = (\mathbb{N}; R_1, \dots, R_k, =, \neq)$ be an abstract structure and $\mathcal{B} = \{B_n\}_{n < \omega}$ be a fixed sequence of subsets of \mathbb{N} .

The enumeration f of the structure \mathfrak{A} is *acceptable with respect to \mathcal{B}* , if for every n ,

$$f^{-1}(B_n) \leq_e f^{-1}(\mathfrak{A})^{(n)} \text{ uniformly in } n.$$

Denote by $\mathcal{E}(\mathfrak{A}, \mathcal{B})$ - the class of all acceptable enumerations.

Definition. The ω - degree spectrum of \mathfrak{A} with respect to $\mathcal{B} = \{B_n\}_{n < \omega}$ is the set

$$DS(\mathfrak{A}, \mathcal{B}) = \{d_e(f^{-1}(\mathfrak{A})) \mid f \in \mathcal{E}(\mathfrak{A}, \mathcal{B})\}$$

Proposition. $DS(\mathfrak{A}, \mathcal{B})$ is upwards closed with respect to total e -degrees.

ω -Co-Spectra

For every $\mathcal{A} \subseteq \mathcal{D}_\omega$ let $\text{co}(\mathcal{A}) = \{\mathbf{b} \mid \mathbf{b} \in \mathcal{D}_\omega \text{ \& } (\forall \mathbf{a} \in \mathcal{A})(\mathbf{b} \leq_\omega \mathbf{a})\}$.

Definition. The ω -co-spectrum of \mathfrak{A} with respect to \mathcal{B} is the set

$$\text{CS}(\mathfrak{A}, \mathcal{B}) = \text{co}(\text{DS}(\mathfrak{A}, \mathcal{B})).$$

ω -Co-Spectra

For every $\mathcal{A} \subseteq \mathcal{D}_\omega$ let $co(\mathcal{A}) = \{\mathbf{b} \mid \mathbf{b} \in \mathcal{D}_\omega \ \& \ (\forall \mathbf{a} \in \mathcal{A})(\mathbf{b} \leq_\omega \mathbf{a})\}$.

Definition. The ω -co-spectrum of \mathfrak{A} with respect to \mathcal{B} is the set

$$CS(\mathfrak{A}, \mathcal{B}) = co(DS(\mathfrak{A}, \mathcal{B})).$$

Proposition. [Selman] For $\mathcal{A} \subseteq \mathcal{D}_e$ we have that $co(\mathcal{A}) = co(\{\mathbf{a} : \mathbf{a} \in \mathcal{A} \ \& \ \mathbf{a} \text{ is total}\})$.

Corollary. $CS(\mathfrak{A}, \mathcal{B}) = co(\{\mathbf{a} \mid \mathbf{a} \in DS(\mathfrak{A}, \mathcal{B}) \ \& \ \mathbf{a} \text{ is a total } e\text{-degree}\})$.

Minimal pair theorem

Theorem. For every structure \mathfrak{A} and every sequence \mathcal{B} there exist total enumeration degrees \mathbf{f} and \mathbf{g} in $DS(\mathfrak{A}, \mathcal{B})$ such that for every ω -enumeration degree \mathbf{a} and $k \in \mathbb{N}$:

$$\mathbf{a} \leq_{\omega} \mathbf{f}^{(k)} \ \& \ \mathbf{a} \leq_{\omega} \mathbf{g}^{(k)} \Rightarrow \mathbf{a} \in CS_k(\mathfrak{A}, \mathcal{B}) .$$

Quasi-Minimal Degree

Theorem. For every structure \mathfrak{A} and every sequence \mathcal{B} , there exists $F \subseteq \mathbb{N}$, such that $\mathbf{q} = d_\omega(F \uparrow \omega)$ and:

- 1 $\mathbf{q} \notin \text{CS}(\mathfrak{A}, \mathcal{B})$;
- 2 If \mathbf{a} is a total e-degree and $\mathbf{a} \geq_\omega \mathbf{q}$ then $\mathbf{a} \in \text{DS}(\mathfrak{A}, \mathcal{B})$
- 3 If \mathbf{a} is a total e-degree and $\mathbf{a} \leq_\omega \mathbf{q}$ then $\mathbf{a} \in \text{CS}(\mathfrak{A}, \mathcal{B})$.

Countable ideals of ω -enumeration degrees

- $I = \text{CS}(\mathfrak{A}, \mathcal{B})$ is a countable ideal.
- $\text{CS}(\mathfrak{A}, \mathcal{B}) = I(\mathbf{f}_\omega) \cap I(\mathbf{g}_\omega)$ where $I(\mathbf{f}_\omega)$ and $I(\mathbf{g}_\omega)$ are the principal ideals of ω -enumeration degrees with greatest elements $\mathbf{f}_\omega = \kappa(\mathbf{f})$ and $\mathbf{g}_\omega = \kappa(\mathbf{g})$, the images of \mathbf{f} and \mathbf{g} under the embedding κ of \mathcal{D}_e in \mathcal{D}_ω .
- Denote by $I^{(k)}$ - the least ideal, containing all k th ω -jumps of the elements of I .

Proposition. [*Ganchev*] $I = I(\mathbf{f}_\omega) \cap I(\mathbf{g}_\omega) \implies I^{(k)} = I(\mathbf{f}_\omega^{(k)}) \cap I(\mathbf{g}_\omega^{(k)})$ for every k .

- $I(\mathbf{f}_\omega^{(k)}) \cap I(\mathbf{g}_\omega^{(k)}) = \text{CS}_k(\mathfrak{A}, \mathcal{B})$ for each k .
- Thus $I^{(k)} = \text{CS}_k(\mathfrak{A}, \mathcal{B})$.

Corollary. $\text{CS}_k(\mathfrak{A}, \mathcal{B})$ is the least ideal containing all k th ω -jumps of the elements of $\text{CS}(\mathfrak{A}, \mathcal{B})$.

Countable ideals of ω -enumeration degrees

There is a countable ideal I of ω -enumeration degrees for which there is no structure \mathfrak{A} and sequence \mathcal{B} such that $I = \text{CS}(\mathfrak{A}, \mathcal{B})$.

- Consider $\mathcal{A} = \{\mathbf{0}_\omega, \mathbf{0}'_\omega, \mathbf{0}''_\omega, \dots, \mathbf{0}_\omega^{(n)}, \dots\}$;
- $I = I(d_\omega(\mathcal{A})) = \{\mathbf{a} \mid \mathbf{a} \in \mathcal{D}_\omega \ \& \ (\exists n)(\mathbf{a} \leq_\omega \mathbf{0}_\omega^{(n)})\}$
- Assume that there is a structure \mathfrak{A} and a sequence \mathcal{B} such that $I = \text{CS}(\mathfrak{A}, \mathcal{B})$
- Then there is a minimal pair \mathbf{f} and \mathbf{g} for $\text{DS}(\mathfrak{A}, \mathcal{B})$, so $I^{(n)} = I(\mathbf{f}_\omega^{(n)}) \cap I(\mathbf{g}_\omega^{(n)})$ for each n .
- But $\mathbf{f}_\omega \geq \mathbf{0}_\omega^{(n)}$ and $\mathbf{g}_\omega \geq \mathbf{0}_\omega^{(n)}$ for each n .
- If $F \in \mathbf{f}$ and $G \in \mathbf{g}$ then $F \geq_T \emptyset^{(n)}$ and $G \geq_T \emptyset^{(n)}$ for every n .
- Then by [Enderton](#) and [Putnam](#) [1970], [Sacks](#) [1971] $F'' \geq_T \emptyset^{(\omega)}$ and $G'' \geq_T \emptyset^{(\omega)}$ and hence $\mathbf{f}'' \geq_T \mathbf{0}_T^{(\omega)}$ and $\mathbf{g}'' \geq_T \mathbf{0}_T^{(\omega)}$.
- Then $\kappa(\iota(\mathbf{0}_T^{(\omega)})) \in I(\mathbf{f}''_\omega) \cap I(\mathbf{g}''_\omega)$.
- But $\kappa(\iota(\mathbf{0}_T^{(\omega)})) \notin I''$ since all elements of I'' are bounded by $\mathbf{0}_\omega^{(k+2)}$ for some k .
- Hence $I'' \neq I(\mathbf{f}''_\omega) \cap I(\mathbf{g}''_\omega)$. A contradiction.

Degree spectra

- Questions:

- ▶ Describe the sets of enumeration degrees which are equal to $DS(\mathfrak{A})$ for some structure \mathfrak{A} .
- ▶ For a countable ideal $I \subseteq \mathcal{D}_\omega$ if there is an exact pair then are there a structure \mathfrak{A} and a sequence \mathcal{B} so that $CS(\mathfrak{A}, \mathcal{B}) = I$?
- ▶ Is it true that for every structure \mathfrak{A} and every sequence \mathcal{B} there exists a structure \mathfrak{B} such that $CS_\omega(\mathfrak{B}) = CS(\mathfrak{A}, \mathcal{B})$? The answer is yes, [Soskov \(2013\)](#), using Marker's extensions